

CHAPTER 6. THE EQUATION $x + y = z$ IN S -INTEGERS.

The results of this chapter have been published in de Weger [1987^a].

6.1. Introduction.

Let S be the set of all positive integers composed of primes from a fixed finite set $\{p_1, \dots, p_s\}$, where $s \geq 3$. This chapter is devoted to the diophantine equation

$$x + y = z \tag{6.1}$$

in $x, y, z \in S$. Without loss of generality we may assume that x, y, z are relatively prime. For any $a \in S$ we define

$$m(a) = \max_{1 \leq i \leq s} \text{ord}_{p_i}(a).$$

It was proved by Mahler [1933] that (6.1) has only finitely many solutions, but his proof is ineffective. An effective version, i.e. an effectively computable upper bound for $m(x \cdot y \cdot z)$ for the solutions x, y, z of (6.1), can be derived from the results of Coates [1969], [1970] and Sprindžuk [1969], since (6.1) can be reduced to a finite number of Thue equations. See also Chapter 1 of Shorey and Tijdeman [1986].

We derive an explicit upper bound in Section 6.2. Section 6.3 is devoted to some details of the p -adic approximation lattices on which the reduction method of Sections 6.4 and 6.5 are based. In Section 6.4 we give a method of solving (6.1) in the one-dimensional case $s = 3$. This method is based on the reduction procedure given in Section 3.10. As an example we find all the solutions of the slightly more general equation $x \pm y = w \cdot z$, where x, y, z are powers of 2, 3 or 5, and $w \in \mathbb{Z}$, $|w| \leq 1000000$, $(w, z) = 1$. In Section 6.5 we give a procedure for solving (6.1) in the multi-dimensional case $s \geq 4$, based on the reduction procedure described in Section 3.11. We work out the example $\{p_1, \dots, p_6\} = \{2, 3, 5, 7, 11, 13\}$, and actually determine all the solutions. This generalizes the result of Alex [1976], who

gave a complete solution of (6.1) for the case $\{ p_1, \dots, p_4 \} = \{ 2, 3, 5, 7 \}$ by elementary arguments. See also Rumsey and Posner [1964] and Brenner and Foster [1982]. We conclude in Section 6.6 with some remarks on the Oesterlé-Masser conjecture, also known as the 'abc'-conjecture, which is related to equation (6.1). In particular, our method of solving (6.1) leads to a method of finding examples that are of interest with respect to the abc-conjecture. Finally, we give tables in Section 6.7.

6.2. Upper bounds.

We give in this section an upper bound for the solutions of (6.1), based on lemma 2.6 (cf. Yu [1987^a]). Note that in our paper de Weger [1987^a] we used the result of van der Poorten [1977] instead of Yu's.

We introduce a lot of notation. Assume that $p_1 < \dots < p_s$. Let q_i be the smallest prime with $q_i \nmid p_i \cdot (p_i - 1)$ for $i = 1, \dots, s$. Put

$$t = [2 \cdot s / 3], \quad P = \prod_{i=1}^s p_i, \quad q = \max_i q_i,$$

$$C_1(2, t) \text{ and } a_1 \text{ as in lemma 2.6 with } n = t,$$

$$U = C_1(2, t) \cdot a_1^t \cdot t^{t+5/2} \cdot q^{2 \cdot t} \cdot (q-1) \cdot \log^2(t \cdot q) \cdot \max_i \frac{(p_i-1) \cdot \left(2 + \frac{1}{p_i-1}\right)^t}{(\log p_i)^{t+2}} \cdot (\log p_s)^t \cdot \left[\log(4 \cdot \log p_s) + \frac{\log p_s}{8 \cdot t} \right],$$

$$C_1 = U / 6 \cdot t, \quad C_2 = U \cdot \log 4,$$

$$V_i = \max(1, \log p_i) \text{ for } i = s-t+1, \dots, s, \quad \Omega = \prod_{i=s-t+1}^s V_i,$$

$$C_3 = 2^{9 \cdot t + 26} \cdot t^{t+4} \cdot \Omega \cdot \log(e \cdot V_{s-1}),$$

$$C_4 = \max \left[7.4, (C_1 \cdot \log(P/p_1) + C_3) / \log p_1 \right],$$

$$C_5 = (C_2 \cdot \log(P/p_1) + C_3 \cdot \log(e \cdot V_s) + 0.327) / \log p_1,$$

$$C_6 = \max \left[C_5, (C_2 \cdot \log(P/p_1) + \log 2) / \log p_1 \right],$$

$$C_7 = 2 \cdot (C_6 + C_4 \cdot \log C_4),$$

$$C_8 = \max \left[p_s, \log \left[2 \cdot (P/p_1)^{p_s} \right] / \log p_1, C_2 + C_1 \cdot \log C_7, C_7 \right] .$$

Now we state the main result.

THEOREM 6.1. *The solutions of (6.1) satisfy $m(x \cdot y \cdot z) \leq C_8$.*

Proof. If we consider instead of (6.1) the equivalent equation

$$x \pm y = z \tag{6.2}$$

then we may assume that $x \cdot y$ has at most t prime divisors, p_{i_1}, \dots, p_{i_t} say. Suppose first that $m(x \cdot y) \leq p_s$. Then

$$p_1^{m(z)} \leq z \leq 2 \cdot \max(x, y) < 2 \cdot (P/p_1)^{p_s} ,$$

hence

$$m(x \cdot y \cdot z) < \max \left[p_s, \log \left[2 \cdot (P/p_1)^{p_s} \right] / \log p_1 \right] \leq C_8 .$$

Next suppose that $m(x \cdot y) \geq p_s$ and $m(z) \geq 2$. Then for some $p = p_{i_1}$,

$$m(z) = \text{ord}_p(z) = \text{ord}_p \left(\frac{x}{y} - 1 \right) = \text{ord}_p \left(\log_p \left(\frac{x}{y} \right) \right) .$$

Put $x/y = \prod_{j=1}^t p_{i_j}^{x_{i_j}}$. Then $m(x \cdot y) = \max_{1 \leq j \leq t} |x_{i_j}|$. We apply Lemma 2.6 (Yu's lemma) with $n = t$, $B_0 = B_n = B' = B = m(x \cdot y)$. Since $m(x \cdot y) \geq p_s$ and $t \geq 2$ we have

$$W = \max \left(\log \left(1 + \frac{3}{4 \cdot t} \cdot B \right), \log B, \log p \right) = \log B .$$

Note that $C_1(p, n)$ is maximal for $p = 2$. We obtain

$$m(z) < C_1 \cdot \log m(x \cdot y) + C_2 . \tag{6.3}$$

Obviously (6.3) is also true if $m(z) < 2$. If in (6.2) the plus sign holds, then

$$(P/p_1)^{m(z)} \geq z > \max(x, y) \geq p_1^{m(x \cdot y)} .$$

By (6.3) and $C_3 > 0$ it then follows that

$$m(x \cdot y) < C_4 \cdot \log m(x \cdot y) + C_6 . \quad (6.4)$$

Next suppose that in (6.2) the minus sign holds. Then we apply Lemma 2.4 to prove (6.4) for this case, as follows. Suppose (6.4) is false. Then

$$\left| \frac{y}{x} - 1 \right| = \frac{z}{x} = \frac{z}{\max(x, y)} \leq \frac{(P/p_1)^{m(z)}}{p_1^{m(x \cdot y)}} < \frac{(P/p_1)^{C_1 \cdot \log m(x \cdot y) + C_2}}{p_1^{C_4 \cdot \log m(x \cdot y) + C_6}} ,$$

which is less than $\frac{1}{2}$, by the definition of C_4 and C_6 . Hence

$$\left| \log \frac{y}{x} \right| < (2 \cdot \log 2) \cdot \left| \frac{y}{x} - 1 \right| < (2 \cdot \log 2) \cdot \frac{(P/p_1)^{C_1 \cdot \log m(x \cdot y) + C_2}}{p_1^{m(x \cdot y)}} .$$

On the other hand, Lemma 2.4 yields

$$\left| \log \frac{y}{x} \right| > \exp \left[-C_3 \cdot (\log m(x \cdot y) + \log(e \cdot V_s)) \right] .$$

Thus we obtain

$$\begin{aligned} m(x \cdot y) \cdot \log p_1 &< \log(2 \cdot \log 2) + (C_1 \cdot \log m(x \cdot y) + C_2) \cdot \log(P/p_1) \\ &+ C_3 \cdot (\log m(x \cdot y) + \log(e \cdot V_s)) \leq (\log p_1) \cdot (C_4 \cdot \log m(x \cdot y) + C_6) . \end{aligned}$$

This contradicts our assumption that (6.4) is false. Consequently (6.4) is true in all cases. Now, by $C_4 > e^2$, Lemma 2.1 yields $m(x \cdot y) < C_7$, and (6.3) then yields $m(x \cdot y \cdot z) < C_8$. \square

Examples. If $s = 3$, $\{ p_1, p_2, p_3 \} = \{ 2, 3, 5 \}$ then $C_8 < 3.98 \times 10^{17}$.
If $s = 6$, $\{ p_1, \dots, p_6 \} = \{ 2, 3, 5, 7, 11, 13 \}$ then $C_8 < 5.60 \times 10^{27}$.

6.3. The p -adic approximation lattices.

As in the proof of Theorem 6.1 we consider (6.2) instead of (6.1). Let p be any of the primes p_1, \dots, p_s . We may assume that $p \nmid x \cdot y$. Rename the other primes as p_0, \dots, p_{s-2} , such that $\text{ord}_p(\log_p(p_0))$ is minimal. For $i = 1, \dots, s-2$ put (cf. Section 3.11)

$$\vartheta_i = -\log_p(p_i) / \log_p(p_0) = \sum_{\ell=0}^{\infty} u_{i,\ell} \cdot p^\ell ,$$

where $u_{i,\ell} \in \{0, 1, \dots, p-1\}$. The ϑ_i take the place of the ϑ'_i of Section 3.11. Then it is clear from Section 3.11 how to define the p-adic approximation lattices Γ_μ for $\mu \in \mathbb{N}_0$. Put

$$\Lambda = \sum_{i=1}^{s-2} x_i \cdot \vartheta_i - x_0.$$

Then Lemma 3.13 yields

$$\begin{aligned} \Gamma_\mu &= \{ (x_1, \dots, x_{s-2}, x_0) \mid |\Lambda|_p \leq p^{-\mu} \} \\ &= \{ (x_1, \dots, x_{s-2}, x_0) \mid \left| \log_p \left(\prod_{i=0}^{s-2} p_i^{x_i} \right) \right|_p \leq p^{-(\mu+\mu_0)} \}, \end{aligned}$$

where $\mu_0 = \text{ord}_p(\log_p(p_0))$. In Section 3.13 we studied the set

$$\Gamma_\mu^* = \{ (x_1, \dots, x_{s-2}, x_0) \mid \left| \prod_{i=0}^{s-2} p_i^{x_i} \pm 1 \right|_p \leq p^{-(\mu+\mu_0)} \},$$

which is a sublattice of Γ_μ . In Lemma 3.17 we showed how a basis of Γ_μ^* can be found from a basis of Γ_μ . In practice this is very easy, especially if for $p \geq 5$ it happens to be possible to choose p_0 such that not only $\text{ord}_p(\log_p(p_0))$ is minimal, but also p_0 is a primitive root (mod p). Then, using the notation of Lemma 3.17 (with \underline{b}_0 as the last element of the basis), choose $\zeta \equiv p_0 \pmod{p}$. Then $k(\underline{b}_i) = 1$, and it follows that $\underline{b}'_i = \underline{b}_i$ for $i = 1, \dots, s-2$. By $\underline{b}_i = (0, \dots, 1, \dots, 0, \vartheta_i^{(\mu)})^T$ we have

$$p_i \cdot p_0^{\vartheta_i^{(\mu)}} \equiv \zeta^{k(\underline{b}_i)} \pmod{p^{\mu+\mu_0}}.$$

If $p_i \equiv p_0^{\alpha_i} \pmod{p}$, then it follows that

$$\begin{aligned} \gamma_i^* &= \alpha_i + \vartheta_i^{(\mu)} \equiv \alpha_i + \sum_{\ell=0}^{\mu-1} u_{i,\ell} \pmod{(p-1)/2} \quad \text{for } i = 1, \dots, s-2, \\ \gamma_0^* &= (p-1)/2. \end{aligned}$$

Lemma 3.14 (with $c_1 = 0, c_2 = 1$) now yields: if

$$\ell(\Gamma_\mu^*) > \surd(s-1) \cdot X_1 \tag{6.5}$$

then (6.2) has no solutions with

$$\mu + \mu_0 \leq \text{ord}_p(z) \leq m(x \cdot y \cdot z) \leq X_1. \tag{6.6}$$

6.4. Reducing the upper bounds in the one-dimensional case.

In Section 3.10 we have described how an upper bound for the solutions of (6.1) in the case $s = 3$ can be reduced. We shall apply that method in this section to the following problem.

THEOREM 6.2. *The diophantine equation*

$$x \pm y = w \cdot z, \tag{6.7}$$

where $x = p_0^{x_0}$, $y = p_1^{x_1}$, $z = p^u$, $(p, p_0, p_1) = (2, 3, 5), (3, 2, 5)$ or $(5, 2, 3)$, $x_0, x_1, u \in \mathbb{N}_0$, $w \in \mathbb{Z}$, $|w| \leq 10^6$, and $p \nmid w$, has exactly 291 solutions for $p = 2$, 412 solutions for $p = 3$, and 570 solutions for $p = 5$. In Table I all solutions with $u \geq 3$ are given. The solutions with $u \leq 2$ satisfy $x_0 \leq 14, x_1 \leq 9$ for $p = 2$, $x_0 \leq 23, x_1 \leq 10$ for $p = 3$, and $x_0 \leq 25, x_1 \leq 15$ for $p = 5$.

Remark. It is easy to find all solutions of (6.7) with $u \leq 2$. The Tables are presented in Section 6.7.

Proof. Put $X = \max_{p=2,3,5} \text{ord}_p(x \cdot y \cdot z)$. The example at the end of Section 6.2 shows that in the case $|w| = 1$ we have $X < 3.98 \times 10^{17}$. It can be checked without difficulties that the effect of the w with $|w| \leq 10^6$ in the proof of Theorem 6.1 can be neglected (it disappears in the rounding off), so that for the solutions of (6.7) also $X < X_0 = 3.98 \times 10^{17}$ holds. Put

$$x/y = p_0^{y_0} \cdot p_1^{y_1}, \quad \vartheta = -\log_p(p_1)/\log_p(p_0).$$

Note that ϑ is a p -adic integer. Define the lattices Γ_μ, Γ_μ^* as in Section 6.3, so Γ_μ is generated by

$$\underline{b}_1 = \begin{bmatrix} 1 \\ \vartheta^{(\mu)} \end{bmatrix}, \quad \underline{b}_0 = \begin{bmatrix} 0 \\ p^\mu \end{bmatrix}.$$

For $p = 2, 3$ we have $\Gamma_\mu^* = \Gamma_\mu$, and for $p = 5$ a basis of Γ_μ^* is

$$\underline{b}_1^* = \underline{b}_1 - \gamma \cdot \underline{b}_0, \quad \underline{b}_0^* = 2 \cdot \underline{b}_0,$$

where $\gamma = 0$ if $\vartheta^{(\mu)}$ is odd, $\gamma = 1$ if $\vartheta^{(\mu)}$ is even. Using the algorithm given in Section 3.10, Fig. 3, we can compute a basis $\underline{c}_1, \underline{c}_2$ of Γ_μ^* that is reduced in the sense that $|\underline{c}_1| = \ell(\Gamma_\mu^*)$. We did so, with μ as

in the following table.

P	P ₀	P ₁	μ ₀	μ	γ	ε ₁ >	u ≤	W	y ₀ ≤	y ₁ ≤
2	3	5	2	143		2.68×10 ²¹	144	10 ⁶ ·2 ¹⁴⁴	114	78
3	2	5	1	91		2.32×10 ²¹	91	10 ⁶ ·3 ⁹¹	182	78
5	2	3	1	65	0	5.28×10 ²²	65	10 ⁶ ·5 ⁶⁵	189	119

The values of $\vartheta(\mu)$ can be found in Table III. Making an exception to our policy, we give the reduced bases of the Γ_{μ}^* below.

$$p = 2 : \quad \varepsilon_1 = \begin{pmatrix} 10 & 00000 & 00100 & 10001 & 10110 & 01110 & 01101 \\ 00001 & 11101 & 00101 & 00100 & 11100 & 01111 & 11010 & 00011 \\ -1 & 00010 & 00110 & 01000 & 01011 & 01110 & 00010 \\ 00101 & 11000 & 00000 & 11100 & 01111 & 01011 & 10111 & 00001 \end{pmatrix},$$

$$\varepsilon_2 = \begin{pmatrix} 10 & 11011 & 10000 & 01011 & 01101 & 11000 & 00111 \\ 11001 & 10100 & 11011 & 00000 & 11111 & 10110 & 10110 & 00001 \\ 10 & 01110 & 11101 & 10111 & 11000 & 00100 & 10101 \\ 00111 & 00001 & 10101 & 00110 & 10011 & 00111 & 00101 & 10101 \end{pmatrix},$$

$$p = 3 : \quad \varepsilon_1 = \begin{pmatrix} -102 & 01121 & 02221 & 00210 & 12120 & 20020 & 22222 & 10212 & 20222 \\ 21002 & 00122 & 21100 & 11102 & 22102 & 20001 & 11222 & 02212 & 21011 \end{pmatrix},$$

$$\varepsilon_2 = \begin{pmatrix} -10 & 12210 & 12111 & 01102 & 02010 & 12112 & 12210 & 21122 & 21011 & 20102 \\ -2 & 22021 & 11012 & 01000 & 12021 & 00211 & 12221 & 22121 & 21220 & 12122 \end{pmatrix},$$

$$p = 5 : \quad \varepsilon_1 = \begin{pmatrix} -211 & 32230 & 21042 & 22023 & 30141 & 33034 & 21420 \\ -22104 & 43102 & 43111 & 03114 & 30134 & 23410 \end{pmatrix},$$

$$\varepsilon_2 = \begin{pmatrix} 340 & 34003 & 02404 & 12120 & 03412 & 22030 & 32211 \\ -414 & 20001 & 42202 & 42210 & 34043 & 20120 & 00432 \end{pmatrix}.$$

From this we found the lower bounds for $|\varepsilon_1|$ given above. They are all larger than $\sqrt{2} \cdot 3.98 \times 10^{17}$. Hence (6.5) holds for $X_1 = X_0$, and then we infer from (6.6) that $u \leq \mu + \mu_0 - 1$, and $|w| \cdot z \leq W$ as shown in the table above. We now find the new upper bounds for $|y_0|$, $|y_1|$ as follows. If in (6.7) the minus sign holds, then, on supposing that $\min(x,y) > W^{10/9}$, we infer

$$|x - y| = |w| \cdot z \leq W < \min(x, y)^{0.9}.$$

By Theorem 5.2(a), the inequality $|x - y| < \min(x, y)^{0.9}$ has no solutions with $\min(x, y) > W$, since $W > 10^{49}$. Hence $\min(x, y) \leq W^{10/9}$, and we infer

$$\max(x, y) \leq \min(x, y) + |w| \cdot z \leq W^{10/9} + W.$$

If in (6.7) the plussign holds, then this inequality follows at once. So now the bounds given in the above table for $|y_0|, |y_1|$ follow from

$$|y_1| \cdot \log p_i \leq \log \max(x, y) \leq \log(W^{10/9} + W).$$

We repeat the procedure with μ as in the following table.

p	μ	γ	$ \underline{c}_1 >$	$\sqrt{2} \cdot X_0 <$	$u \leq$	W	$ y_0 \leq$	$ y_1 \leq$
2	16		167.7	161.3	17	$10^6 \cdot 2^{17}$	31	21
3	13		535.8	257.4	13	$10^6 \cdot 3^{13}$	49	21
5	7	1	276.1	267.3	7	$10^6 \cdot 5^7$	49	31

The numbers are now so small that the computations can be performed by hand. For example, for $p = 5$, the lattice Γ_7^* is generated by

$$\underline{b}_1^* = \begin{bmatrix} 1 \\ -45607 \end{bmatrix}, \quad \underline{b}_0^* = \begin{bmatrix} 0 \\ 156250 \end{bmatrix},$$

and a reduced basis is

$$\underline{c}_1 = \begin{bmatrix} 185 \\ 205 \end{bmatrix}, \quad \underline{c}_0 = \begin{bmatrix} -394 \\ 408 \end{bmatrix}.$$

We find upper bounds for u and W as given in the above table. In all three cases, $W^{10/9} < 10^{15}$. On supposing $\min(x, y) > 10^{15}$ we infer

$$|x - y| = |w| \cdot z \leq W < 10^{15 \cdot 0.9} \leq \min(x, y)^{0.9}.$$

By Theorem 5.2(a) we see that the inequality $|x - y| < \min(x, y)^{0.9}$ has only two solutions: $(x, y) = (2^{65}, 5^{28}), (2^{84}, 3^{53})$. However, both have $|x - y| > 10^{15 \cdot 0.9}$. So we infer $\min(x, y) \leq 10^{15}$, hence by $\max(x, y) \leq 10^{15} + W$ we obtain the bounds for $|y_0|, |y_1|$ as given above. These bounds are small enough to admit enumeration of the remaining cases. \square

Remark. The computer calculations for the above proof took less than 1 sec.

6.5. Reducing the upper bounds in the multi-dimensional case.

In Section 3.11 we have described how an upper bound for the solutions of (6.1) in the case $s \geq 3$ can be reduced. We shall apply that method in this section to the following problem.

THEOREM 6.3. *The diophantine equation*

$$x + y = z \tag{6.8}$$

in $x, y, z \in S = \{ 2^{x_1} \dots 13^{x_6} \mid x_i \in \mathbb{N}_0 \text{ for } i = 1, \dots, 6 \}$ with $(x,y) = 1$ and $x \leq y$ has exactly 545 solutions. Of them, 514 satisfy

$$\begin{aligned} \text{ord}_2(x \cdot y \cdot z) \leq 12, \quad \text{ord}_3(x \cdot y \cdot z) \leq 7, \quad \text{ord}_5(x \cdot y \cdot z) \leq 5, \\ \text{ord}_7(x \cdot y \cdot z) \leq 4, \quad \text{ord}_{11}(x \cdot y \cdot z) \leq 3, \quad \text{ord}_{13}(x \cdot y \cdot z) \leq 4. \end{aligned}$$

The remaining 31 solutions are given in Table II.

Remark. From Theorem 6.3 it is not much effort to find all 545 solutions of (6.8).

Proof. In the example at the end of Section 6.2 we have seen that $m(x \cdot y \cdot z) < X_0 = 5.60 \times 10^{27}$. With the notation of Section 6.3 we choose the following parameters.

p	p_0	p_1	p_2	p_3	p_4	μ_0	μ	γ_0^*	γ_1^*	γ_2^*	γ_3^*	γ_4^*
2	3	5	7	11	13	2	605					
3	2	5	7	11	13	1	385					
5	2	3	7	11	13	1	275	2	0	1	1	1
7	3	2	5	11	13	1	220	3	0	1	1	0
11	2	3	5	7	13	1	165	5	-2	0	1	1
13	2	3	5	7	11	1	165	6	2	1	2	3

We computed the six values of the $\vartheta_i^{(\mu)}$ for $i = 1, 2, 3, 4$ (and give them in Table III), and the reduced bases of the six lattices Γ_μ^* , by the L^3 -algorithm. Thus we obtained:

p	$\ell(\Gamma_\mu^*) \geq \underline{\varepsilon}_1 /4 >$	$\text{ord}_p(x \cdot y \cdot z) \leq$
2	4.70×10^{35}	606
3	1.15×10^{36}	385
5	6.27×10^{37}	275
7	3.17×10^{36}	220
11	5.74×10^{33}	165
13	1.73×10^{36}	165

These lower bounds for $\ell(\Gamma_\mu^*)$ are all larger than $\sqrt{5} \cdot 5.60 \times 10^{27}$ (note that we have a very large margin here, we could have taken the μ 's probably about 20% smaller). So we apply Lemma 3.14 for $X_1 = X_0 = 5.60 \times 10^{27}$. For every p we thus find $\text{ord}_p(z) \leq \mu + \mu_0$. Since equation (6.2) is invariant under permutations of x, y, z, we even have $\text{ord}_p(x \cdot y \cdot z) \leq \mu + \mu_0$, as shown in the above table. Hence $m(x \cdot y \cdot z) \leq 606$.

We repeated the procedure with $X_0 = 606$ and μ as in the following table. After computing the reduced bases of the six lattices Γ_μ^* we found the following data. Note that in all cases $\ell(\Gamma_\mu^*) \geq \sqrt{5} \cdot 606$.

p	μ	γ_0^*	γ_1^*	γ_2^*	γ_3^*	γ_4^*	$\ell(\Gamma_\mu^*) >$	$\text{ord}_p(x \cdot y \cdot z) \leq$
2	66						1909	67
3	42						2304	42
5	30	2	0	0	1	1	3417	30
7	24	3	1	0	-1	1	2391	24
11	18	5	0	2	-2	1	1443	18
13	18	6	0	-1	-1	2	3196	18

Hence $m(x \cdot y \cdot z) \leq 67$. Next, we repeated the procedure with $X_0 = 67$, and μ as in the following table. We found

p	μ	γ_0^*	γ_1^*	γ_2^*	γ_3^*	γ_4^*	$\ell(\Gamma_\mu^*) >$	$\text{ord}_p(x \cdot y \cdot z) \leq$
2	55						364	56
3	35						301	35
5	25	2	1	1	1	0	622	25
7	20	3	1	-1	1	0	693	20
11	15	5	1	2	-2	-2	192	15
13	15	6	1	0	3	2	658	15

Hence $m(x \cdot y \cdot z) \leq 56$.

To find the solutions of (6.2) with $\text{ord}_p(x \cdot y \cdot z)$ below the bounds given in the above table, we followed the following procedure. Suppose that we are at a certain moment interested in finding the solutions with $\text{ord}_p(x \cdot y \cdot z) \leq f(p)$ where $f(p)$ is given for $p = 2, \dots, 13$. Choose p , and $\mu < f(p) - \mu_0$, and consider the lattice Γ_μ^* for these values of p, μ . If a solution x, y, z of (6.2) exists with $\text{ord}_p(z) \geq \mu + \mu_0$, then the vector $(x_1, \dots, x_4, x_0)^T$ with $x_i = \text{ord}_{p_i}(x/y)$ for $i = 0, \dots, 4$, is in the lattice. Its length is bounded by $\sqrt{f(p_0)^2 + \dots + f(p_4)^2}$. All vectors in Γ_μ^* with length below this bound can be computed by the algorithm of Fincke and Pohst, as given in Section 3.6. Then all solutions of (6.2) corresponding to lattice points can be selected. Then we replace $f(p)$ by $\mu + \mu_0 - 1$, and we repeat the procedure for newly chosen p, μ .

We performed this procedure, starting with the bounds for $\text{ord}_p(x \cdot y \cdot z)$ given in the above table for $f(p)$, and with p, m as in the table on the next page. Here, # stands for the number of solutions of (5.2) found at that stage. At the end we have $f(2) = 4$, $f(p) = 1$ for $p = 3, \dots, 13$. The remaining solutions can be found by hand. \square

Remark. Theorems 6.2 and 6.3 have applications in group theory (cf. Alex [1976]). We use Theorem 6.3 in Section 7.2.

Remark. The computer calculations for the proof of Theorem 6.3 took 438 sec., of which 412 were used for the first reduction step. In this first step we applied the L^3 -algorithm in 11 steps (cf. Section 3.5), which cost on average about 60 sec. per lattice. The remaining 50 sec. were mainly used for the computation of the 24 $\vartheta_i^{(\mu)}$'s.

6.6. Examples related to the abc-conjecture.

Let x, y, z be positive integers. Put

$$G = \prod_{\substack{p|xyz \\ p \text{ prime}}} p.$$

For all x, y, z with $(x, y) = 1$ and $x + y = z$ we define

$$c(x, y, z) = \log z / \log G.$$

Recently, Oesterlé posed the problem to decide whether there exists an

p	m	#	p	m	#	p	m	#
2	44	-	2	13	1	2	10	2
3	28	-	2	12	2	2	9	3
5	20	-	2	11	2	2	8	6
7	16	-	3	13	-	2	7	15
11	12	-	3	12	-	2	6	16
13	12	-	3	11	-	2	5	26
2	33	-	3	10	1	2	4	31
3	21	-	3	9	1	2	3	44
5	15	-	3	8	1	3	6	5
7	12	-	3	7	6	3	5	8
11	9	-	5	9	-	3	4	16
13	9	-	5	8	-	3	3	35
2	22	-	5	7	-	3	2	54
3	14	-	5	6	-	3	1	87
5	10	-	5	5	6	5	4	1
7	8	-	7	7	-	5	3	5
11	6	-	7	6	-	5	2	18
13	6	-	7	5	1	5	1	36
2	21	-	7	4	4	7	3	-
2	20	-	11	5	-	7	2	6
2	19	-	11	4	1	7	1	18
2	18	-	11	3	4	11	2	1
2	17	-	13	5	-	11	1	8
2	16	-	13	4	-	13	2	-
2	15	-	13	3	1	13	1	4
2	14	-						

absolute constant C such that $c(x,y,z) < C$ for all x, y, z . Masser conjectured the stronger assertion that $c(x,y,z) < 1 + \epsilon$, when z exceeds some bound depending on ϵ only, for all $\epsilon > 0$. For a survey of related results and conjectures, see Stewart and Tijdeman [1986] and Vojta [1987].

It might be interesting to have some empirical results on $c(x,y,z)$, and to search for x, y, z for which it is large. From the preceding sections it may be clear that such x, y, z correspond to relatively short vectors in appropriate p -adic approximation lattices.

As a byproduct of the proofs of Theorems 5.5 and 6.3 we computed the value of $c(x,y,z)$, corresponding to many short vectors that we came across in performing the algorithm of Fincke and Pohst. All examples that we found with $c(x,y,z) \geq 1.4$ are listed below. Our search was rather unsystematic, so we do not guarantee that this list is complete in any sense. The largest value for $c(x,y,z)$ that occurred is 1.626 , which was reached by

$$x = 11^2 = 121, y = 3^2 \cdot 5^6 \cdot 7^3 = 48234375, z = 2^{21} \cdot 23 = 48234496 .$$

This example was found on September 20, 1985, and has not yet been beaten, to the author's knowledge.

x	y	z	$c(x,y,z)$
11^2	$3^2 \cdot 5^6 \cdot 7^3$	$2^{21} \cdot 23$	1.62599
1	$2 \cdot 3^7$	$5^4 \cdot 7$	1.56789
7^3	3^{10}	$2^{11} \cdot 29$	1.54708
$5^2 \cdot 7937$	7^{13}	$2^{18} \cdot 3^7 \cdot 13^2$	1.49762
11^2	$3^9 \cdot 13$	$2^{11} \cdot 5^3$	1.48887
37	2^{15}	$3^8 \cdot 5$	1.48291
$2^7 \cdot 5^2$	$7^6 \cdot 41$	13^6	1.46192
1	$2^5 \cdot 3 \cdot 5^2$	7^4	1.45567
$2^{19} \cdot 13 \cdot 103$	7^{11}	$3^{11} \cdot 5^3 \cdot 11^2$	1.45261
1	$2^{12} \cdot 5^3$	$3^5 \cdot 7^2 \cdot 43$	1.44331
1	$2^4 \cdot 3^7 \cdot 547$	$5^8 \cdot 7^2$	1.43906
$2^{10} \cdot 7$	5^7	$3^8 \cdot 13$	1.43501
3	5^3	2^7	1.42657
5	3^{11}	$2^{10} \cdot 173$	1.41268

These results do not seem to yield any heuristical evidence for the truth or falsity of the abc-conjecture.

6.7. Tables.

Table I. (Theorem 6.2.)

$$p = 2, p_0 = 3, p_1 = 5$$

N_0	$p_0^{N_0}$	N_1	$p_1^{N_1}$	sign	u	w
2	9	10	9765625	-1	4	-610351
10	59049	10	9765625	-1	4	-606661
4	81	12	244140625	-1	9	-476837
6	729	10	9765625	-1	5	-305153
2	9	8	390625	-1	3	-48827
6	729	8	390625	-1	3	-48737
10	59049	8	390625	-1	3	-41447
14	4782969	10	9765625	-1	7	-38927
4	81	8	390625	-1	4	-24409
0	1	8	390625	-1	5	-12207
8	6561	8	390625	-1	6	-6001
0	1	6	15625	-1	3	-1953
4	81	6	15625	-1	3	-1943
8	6561	6	15625	-1	3	-1133
6	729	6	15625	-1	4	-931
2	9	4	625	-1	3	-77
2	9	6	15625	-1	8	-61
0	1	4	625	-1	4	-39
4	81	4	625	-1	5	-17
0	1	2	25	-1	3	-3
2	9	2	25	-1	4	-1
1	3	1	5	1	3	1
1	3	3	125	1	7	1
2	9	0	1	-1	3	1
3	27	1	5	1	5	1
4	81	0	1	-1	4	5
4	81	2	25	-1	3	7
6	729	2	25	-1	6	11
6	729	4	625	-1	3	13
3	27	3	125	1	3	19
5	243	3	125	1	4	23
5	243	1	5	1	3	31
7	2187	5	3125	1	6	83
6	729	0	1	-1	3	91
7	2187	1	5	1	4	137
11	177147	1	5	1	10	173
3	27	5	3125	1	4	197
8	6561	0	1	-1	5	205
7	2187	3	125	1	3	289
8	6561	4	625	-1	4	371

Table continued

Table I. (cont.)

x_0	$p_0^{x_0}$	x_1	$p_1^{x_1}$	sign	u	w
1	3	5	3125	1	3	391
5	243	5	3125	1	3	421
9	19683	3	125	1	5	619
8	6561	2	25	-1	3	817
10	59049	6	15625	-1	5	1357
5	243	7	78125	1	5	2449
9	19683	1	5	1	3	2461
9	19683	5	3125	1	3	2851
10	59049	2	25	-1	4	3689
12	531441	4	625	-1	7	4147
1	3	7	78125	1	4	4883
9	19683	7	78125	1	4	6113
13	1594323	7	78125	1	8	6533
10	59049	4	625	-1	3	7303
10	59049	0	1	-1	3	7381
12	531441	8	390625	-1	4	8801
3	27	7	78125	1	3	9769
7	2187	7	78125	1	3	10039
11	177147	5	3125	1	4	11267
3	27	9	1953125	1	7	15259
11	177147	3	125	1	3	22159
11	177147	7	78125	1	3	31909
12	531441	0	1	-1	4	33215
12	531441	6	15625	-1	3	64477
12	531441	2	25	-1	3	66427
11	177147	9	1953125	1	5	66571
13	1594323	3	125	1	4	99653
7	2187	9	1953125	1	4	122207
14	4782969	2	25	-1	5	149467
13	1594323	1	5	1	3	199291
13	1594323	5	3125	1	3	199681
1	3	9	1953125	1	3	244141
5	243	9	1953125	1	3	244171
9	19683	9	1953125	1	3	246601
14	4782969	6	15625	-1	4	297959
13	1594323	9	1953125	1	3	443431
15	14348907	5	3125	1	5	448501
14	4782969	8	390625	-1	3	549043
14	4782969	4	625	-1	3	597793
14	4782969	0	1	-1	3	597871
16	43046721	0	1	-1	6	672605
9	19683	11	48828125	1	6	763247
15	14348907	1	5	1	4	896807

Table continued

Table I. (cont.)

$p = 3, p_0 = 2, p_1 = 5$

x_0	$p_0^{x_0}$	x_1	$p_1^{x_1}$	sign	u	w
14	16384	10	9765625	-1	4	-120361
9	512	9	1953125	-1	3	-72319
4	16	8	390625	-1	3	-14467
12	4096	6	15625	-1	3	-427
7	128	5	3125	-1	4	-37
2	4	4	625	-1	3	-23
1	2	2	25	1	3	1
5	32	1	5	-1	3	1
6	64	3	125	1	3	7
11	2048	4	625	1	5	11
9	512	0	1	1	3	19
10	1024	2	25	-1	3	37
3	8	6	15625	1	4	193
15	32768	3	125	-1	4	403
14	16384	1	5	1	3	607
17	131072	7	78125	-1	3	1961
16	65536	5	3125	1	3	2543
8	256	7	78125	1	3	2903
19	524288	2	25	1	4	6473
18	262144	0	1	-1	3	9709
23	8388608	1	5	-1	6	11507
13	8192	8	390625	1	3	14771
22	4194304	8	390625	-1	5	15653
10	1024	11	48828125	1	7	22327
18	262144	9	1953125	1	4	27349
20	1048576	4	625	-1	3	38813
0	1	9	1953125	1	3	72338
21	2097152	6	15625	1	3	78251
5	32	10	9765625	1	3	361691
24	16777216	3	125	1	3	621383
23	8388608	10	9765625	1	3	672379
26	67108864	7	78125	1	4	829469

$p = 5, p_0 = 2, p_1 = 3$

x_0	$p_0^{x_0}$	x_1	$p_1^{x_1}$	sign	u	w
12	4096	16	43046721	-1	3	-344341
5	32	15	14348907	-1	3	-114791
7	128	1	3	-1	3	1
6	64	8	6561	1	3	53
14	16384	2	9	-1	3	131
13	8192	9	19683	1	3	223
20	1048576	10	59049	1	3	8861
21	2097152	3	27	-1	3	16777

Table II. (Theorem 6.3.)

X	Y	z	ord _p (x)							ord _p (y)							ord _p (z)						
			p=2	3	5	7	11	13	p=2	3	5	7	11	13	p=2	3	5	7	11	13			
2401	4160	6561	0	0	0	4	0	0	6	0	1	0	0	1	0	8	0	0	0	0			
875	6561	7436	0	0	3	1	0	0	0	8	0	0	0	0	2	0	0	0	1	2			
1183	6561	7744	0	0	0	1	0	2	0	8	0	0	0	0	6	0	0	0	2	0			
1125	8192	9317	0	2	3	0	0	0	13	0	0	0	0	0	0	0	0	1	3	0			
1183	8192	9375	0	0	0	1	0	2	13	0	0	0	0	0	0	1	5	0	0	0			
16	14625	14641	4	0	0	0	0	0	0	2	3	0	0	1	0	0	0	0	4	0			
81	14560	14641	0	4	0	0	0	0	5	0	1	1	0	1	0	0	0	0	4	0			
1936	13689	15625	4	0	0	0	2	0	4	0	0	0	2	0	0	0	6	0	0	0			
3718	11907	15625	1	0	0	0	1	2	0	5	0	2	0	0	0	0	6	0	0	0			
5824	9801	15625	6	0	0	1	0	1	0	4	0	0	2	0	0	0	6	0	0	0			
49	16335	16384	0	0	0	2	0	0	0	3	1	0	2	0	14	0	0	0	0	0			
2695	13689	16384	0	0	1	2	1	0	0	4	0	0	0	2	14	0	0	0	0	0			
8019	8788	16807	0	6	0	0	1	0	2	0	0	0	0	3	0	0	0	5	0	0			
3584	14641	18225	9	0	0	1	0	0	0	0	0	0	4	0	0	6	2	0	0	0			
1625	16807	18432	0	0	3	0	0	1	0	0	0	0	5	0	0	2	0	0	0	0			
3993	16807	20800	0	1	0	0	3	0	0	0	0	0	5	0	0	0	2	0	0	1			
49	28512	28561	0	0	0	2	0	0	5	4	0	0	1	0	0	0	0	0	0	4			
12936	15625	28561	3	1	0	2	1	0	0	0	6	0	0	0	0	0	0	0	0	4			
22000	6561	28561	4	0	3	0	1	0	0	8	0	0	0	0	0	0	0	0	0	4			
15625	17303	32928	0	0	6	0	0	0	0	0	0	0	3	1	5	1	0	3	0	0			
507	32768	33275	0	1	0	0	0	2	15	0	0	0	0	0	0	0	2	0	3	0			
10985	41503	52488	0	0	1	0	0	3	0	0	0	3	2	0	3	8	0	0	0	0			
10000	49049	59049	4	0	4	0	0	0	0	0	0	3	1	1	0	10	0	0	0	0			
14641	46875	61516	0	0	0	0	4	0	0	1	6	0	0	0	2	0	0	1	0	3			
7168	78125	85293	10	0	0	1	0	0	0	0	7	0	0	0	0	8	0	0	0	1			
20449	97200	117649	0	0	0	0	2	2	4	5	2	0	0	0	0	0	0	6	0	0			
13	151250	151263	0	0	0	0	0	1	1	0	4	0	2	0	0	2	0	5	0	0			
12005	161051	173056	0	0	1	4	0	0	0	0	0	0	5	0	10	0	0	0	2	0			
121	255879	256000	0	0	0	0	2	0	0	9	0	0	0	1	11	0	3	0	0	0			
2197	583443	585640	0	0	0	0	0	3	0	5	0	4	0	0	3	0	1	0	4	0			
91	1771470	1771561	0	0	0	1	0	1	1	11	1	0	0	0	0	0	0	0	6	0			

Table III.

- $\log_2 5 / \log_2 3 =$	0.10101 11101 00001 11110 11000 10101 00000 01001 11101 10000 01011 11100 00001 11010 00000 00001
	00010 11100 11100 10111 01001 01110 11000 01010 01110 01010 11110 00000 01011 01111 01110 01010 11101
	10010 01001 00001 10100 00111 00001 11111 01111 00001 01101 01100 00010 01101 11100 01101 11011 01011
	00001 00111 11011 11001 01000 10001 10010 01011 00000 01100 01111 01101 00110 11110 00000 01100 11010
	00000 11101 10101 11100 10010 11101 01011 10001 01100 00110 01001 01000 00110 01100 11101 11101 10100
	00110 00011 01111....
- $\log_2 7 / \log_2 3 =$	0.01001 01011 01111 11100 11010 01111 11111 10010 01000 11100 00011 00100 01011 11010 10001 00000 01110
	01011 00101 10010 00111 10111 01001 10001 11000 11011 01011 01001 11000 11011 10010 10001 00001 01110
	11000 01001 01110 10000 01101 11000 11001 00011 00011 01100 11110 00011 01100 00011 10110 00001 10111
	10010 01000 00011 01011 10010 11001 10000 01101 10111 01001 01000 01101 00011 11001 10000 00011 00011
	11100 01110 11110 10101 00101 01110 11100 10000 01011 00100 01100 11100 00100 01110 10001 00001 10111
	10111 01100 10110 00111 00101 10011 00101 01011 00101 01011 11100 11011 01011 10110 10001 10001 10111
	10001 11110 1010....
- $\log_2 11 / \log_2 3 =$	0.10011 01110 00001 01001 00110 01010 01110 00100 00101 10000 01000 11000 01001 10110 00011 11101
	10100 01101 01101 10111 10110 01100 10110 01100 10110 00000 11000 11010 01011 10010 00001 01110 10100
	01100 10100 01000 10101 00010 01011 10111 10000 00001 01000 11010 10110 00001 10110 11110 10010 10010
	11101 11100 10010 00100 11000 00000 01110 01000 00001 01100 10010 00101 10010 01111 10000 10101 01000
	10001 11001 10011 11110 10000 11001 10110 10000 11001 10110 10000 11001 10110 10000 10110 11110 10010
	10101 01100 11011 00111 11000 00111 11100 00111 11100 01101 11100 01101 10110 10000 10110 11110 10010
	10101 11001 1111....
- $\log_2 13 / \log_2 3 =$	0.11011 10110 10100 10001 01100 01111 10001 00110 00001 01110 01011 10110 11100 00111 11111 11001 11110
	00011 11110 01000 10010 11011 11101 01000 11101 11101 00001 10101 00000 00001 01101 10010 00001 10101
	01011 11100 10011 11011 00000 01110 10100 00011 10000 00011 10101 00001 00000 10010 10010 10010 00100
	10011 00001 01000 10101 10110 01001 10110 01001 11100 00001 11100 10001 01110 10001 01111 01111 00001
	10011 00000 01011 10101 10111 11100 01011 01011 11000 10001 11110 00011 10110 11011 10110 10100 01110
	01000 00111 00000 00000 00011 10111 10100 01111 10010 01100 00011 00111 10110 11010 10110 10111 10001
	00000 01100 1001....

Table III. (cont.)

$-\log_3 5 / \log_3 2 =$	0.11022	12121	22001	12010	21102	10210	10022	20212	20010	10112	22201	21021	21022	10000	22620	12012	02022	21001	00012	02020
	21210	12202	12200	00099	10120	00211	12021	10120	02109	10222	22122	01201	21111	11121	11001	20222	10000	20121	22221	01002
	20220	12211	22211	00100	20202	00012	11112	10122	21001	21200	12201	12220	11100	01102	20010	11102	10222	00020	21202	21112
	20201	21100	11212	22222	21120	02020	12121	02122	11111	10001	10220	21622	10012	11212	20001	10211	02120	02122	1....	
$-\log_3 7 / \log_3 2 =$	0.20101	10202	20011	12121	01102	11100	01210	20120	02122	02012	20202	00121	21200	01201	11120	11211	11212	22100	00100	22201
	20021	11112	00122	00011	22130	00000	22011	11100	12010	22110	12122	00222	10220	21102	20001	02101	00121	11002	11012	12201
	21011	20100	01110	02000	21222	12010	02201	22612	01022	02021	00210	10221	00221	20202	02222	22122	00100	12021	21220	02220
	20000	00002	00111	11221	11002	20102	12212	12012	22122	00211	01210	01102	21010	20121	01020	11111	20002	10122	2....	
$-\log_3 11 / \log_3 2 =$	0.21112	20101	00222	20222	01212	01100	12100	01201	01111	01212	01210	20121	20051	12021	01122	21202	12020	00212	11102	11002
	01001	10200	22202	02001	20022	10221	00010	10011	22220	01021	02121	00211	22210	21101	22012	11111	02010	00221	00102	20111
	20202	01201	01220	22022	11221	10121	10202	10011	11002	10220	22110	21121	00112	02122	21200	01021	21002	21002	10010	00110
	00101	12202	12000	21012	11010	11020	00222	00012	11201	11010	00122	01120	22200	20112	12122	10202	01211	00210	2....	
$-\log_3 13 / \log_3 2 =$	0.10221	02211	12122	22010	10002	01221	00121	02020	11201	02021	02112	21010	20122	02001	02112	21012	10222	01002	01200	01211
	10111	21100	12121	11010	02000	02212	11111	21220	22020	02000	01222	12112	02100	10110	20002	10222	02112	20112	11100	00211
	20012	11102	22220	00112	00001	11110	11102	22201	01122	22211	22201	11011	22201	01200	22121	02101	22222	22002	01010	01021
	12020	20111	12102	00011	02002	02000	10211	00222	12202	02202	20212	22012	01222	20220	11211	20021	11111	00000	2....	
$-\log_5 3 / \log_5 2 =$	0.33002	02003	04011	23120	04012	01011	00004	43204	30304	00023	14333	12413	43420	40302	10202	44104	32433	26432	03021	12311
	33044	40231	04112	33230	00242	14232	14400	31104	42112	44033	11014	44344	12114	44211	32120	43131	34041	00411	34233	41410
	24120	42032	43014	21421	40044	01142	21004	42021	14011	10404	00214	51110	04441	42431	24423	02433			
$-\log_5 7 / \log_5 2 =$	0.03044	34433	10114	43203	12033	14002	12341	31312	03421	00343	41423	00040	24241	22103	14260	32214	11401	42230	13040	33404
	04310	43034	13233	23241	43062	44411	41124	22443	42612	30420	11223	43101	01000	42112	10643	34210	03410	14414	02220	24443
	13332	33123	23331	20323	44440	13210	14403	32122	03040	31123	04212	22443	44223	23133	02003	12440			

Table III. (cont.)

- log ₅ 11 / log ₂ 2 =	21041	12112	42420	00220	41143	12040	32144	21100	01304	24013	43401	23313	12022	34404
	0.44032	21012	13124	21134	03320	33622	21041	12112	42420	00220	41143	12040	32144	21100
		12413	10214	30123	11014	24110	42444	42030	02413	20241	22304	23423	13414	03234
		31022	33142	14441	44113	21413	23132	31413	32032	01221	40210	24101	30133	13110
														13400
														22110
														23334....
														00330
														01104
														44410
														44113
- log ₅ 13 / log ₂ 2 =	32420	14114	41224	04403	11334	43213	33303	03130	32244	11153	43543	23422	11320	41041
	0.12423	02224	01323	24314	23021	32420	14114	41224	04403	11334	43213	33303	03130	32244
		61320	34110	03024	40012	23213	10014	41441	04420	40114	00021	33224	30103	03243
		22012	30332	22422	43110	13302	34431	13241	13230	44204	14432	33210	24121	13144
														03230
														14301
														3040....
														00013
														04203
														43134
- log ₇ 2 / log ₃ 2 =	40631	52354	45764	60036	13315	13044	46363	40432	02366	04135	21304	53356	32205	44546
	0.20603	14521	11264	52354	45764	60036	13315	13044	46363	40432	02366	04135	21304	53356
		40631	52356	64053	50031	20136	46625	03465	02235	11551	46123	25164	52364	25520
		41326	32413	65633	52502	526....								12240
														64220
														00164
														43634
														02066
														41264
														61233
- log ₇ 5 / log ₃ 2 =	0.62250	35002	24045	66544	01041	43506	34535	04453	26545	45453	33261	65353	53330	22443
		42366	51054	13301	06465	43020	41555	41121	64255	11350	55053	64515	44465	36222
		21343	24510	62633	00155	361....								25605
														66346
														16142
														31340
														45522
														31033
														14255
- log ₇ 11 / log ₃ 2 =	0.25035	56505	33552	02331	10224	32143	50543	02561	42352	23430	26326	53466	23462	31210
		00441	56630	66142	56540	44042	52255	15314	10131	40626	10545	02504	04254	45066
		41366	64215	64014	56645	550....								50232
														65102
														41255
														41254
														54156
														34054
- log ₇ 13 / log ₃ 2 =	0.21305	11055	56501	22565	55610	32506	13150	66465	56420	46445	21450	16426	11613	41010
		43655	56120	34610	53642	61544	36122	61225	42410	23035	15004	26220	14444	23632
		24255	36603	45452	66563	536....								33426
														15605
														51104
														04520
														65502
														65542

Table III. (cont.)

- log ₁₁ 3 / log ₁₁ 2 =																			
0.08248	A4245	06166	43468	58202	44A56	73171	16753	A203A	8A543	28431	86731	11411	4A296	993A7	31A79	00421	95444	80670	57433
59439	78064	34745	1A710	64682	08044	81761	27049	03452	3661A	40979	29601	898A4	50.....						
- log ₁₁ 5 / log ₁₁ 2 =																			
0.351A9	7223A	31378	09193	42445	306A3	96588	11862	48667	AA6A2	39A03	77139	01693	21678	33652	12687	95AA8	24190	78276	28711
08399	68022	2A607	55A17	2231A	80798	76947	73936	21855	30A1A	95324	1A8A3	82999	67.....						
- log ₁₁ 7 / log ₁₁ 2 =																			
0.44804	92167	71327	83472	37453	00781	3256A	2A367	85671	88907	799A1	4AAAA	784A1	29329	A6950	17481	86846	17379	94130	77091
29354	33161	9A146	03746	52A14	20214	22541	58A91	50337	7954A	89A01	43809	A8152	52.....						
- log ₁₁ 13 / log ₁₁ 2 =																			
0.9011A	94962	52990	39096	3A68A	7556A	1A4A3	44758	57692	20188	42770	072A3	9A977	8819A	97518	14396	07360	899A2	99391	26176
84077	81181	54473	58532	58A01	91643	28056	63940	99265	27989	37450	85913	91289	56.....						
- log ₁₃ 3 / log ₁₃ 2 =																			
0.621B3	15581	0A077	3B5C8	49202	39A32	82105	848C7	70988	B863B	75151	52114	5C25A	04902	6B6C7	377B9	3122B	5CAC0	13945	A2471
9B4BA	A79C2	7A91C	5A989	C392A	CC16A	A20A1	75C6B	06BB6	8A3D9	C782B	AA70C	AB218	C9.....						
- log ₁₃ 5 / log ₁₃ 2 =																			
0.44570	79C51	73665	3796C	B7C61	335A0	79906	2B429	51211	4900B	481B1	621AB	2AC77	C2291	1662A	BB03A	8CB9C	77331	74992	11C07
BB101	10301	77310	B8B28	83AB2	57975	7C697	57928	23B72	297CB	0A414	32B3C	A67A8	48.....						
- log ₁₃ 7 / log ₁₃ 2 =																			
0.11C78	9C71A	63110	51424	42CA9	0AAAT	B225B	B0281	501B1	976C2	3C05B	09CA3	AB8E3	C3251	838AC	72502	A1844	03603	644A8	A8501
175BB	BB1C1	30466	223C6	C98B4	564C2	47140	28856	C8676	15C50	12892	A3317	163C8	CA.....						
- log ₁₃ 11 / log ₁₃ 2 =																			
0.1760A	A080C	20874	BB876	B2162	75989	CB19B	B7CC2	26RB7	87093	5A833	A9375	AB4BA	8C0BC	1A698	96C6B	A9411	34B75	4B718	63BC3
571A9	14566	8619B	A4B95	B4244	452A8	29623	49AA5	CB804	AC61A	CC513	08855	79185	43.....						